## Game-Theoretic Models for Rapid Operational Airlift Network Design in Contested Environments



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**Question:** How should AMC allocate airlift capacity in a contested environment?



# **Measuring Airlift Capacity**

#### Definition

The US Air Force uses "Maximum on Ground (MOG)" as a measure of an airfield's capacity for processing aircraft.

Reference: AFPAM 10-1403

- Indicates, e.g., the amount of support equipment and personnel assigned to an airfield.
- ▶ 1 MOG  $\approx$  8 C-17 equivalents per day



A C-17 engine being loaded onto a C-17 at Yokota Air Base, Japan. (Source: Air Mobility Command)

# The Problem, and Our Contributions

**The Problem:** Determine how much MOG should be at each airfield in the airlift network.

**Importance:** Airlift is an important means of supplying distributed forces over long distances.

Solutions Without Our Help: Recent work at the Naval Postgraduate School [Whitlow 2022] developed an optimization model of the problem, called the Airlift Optimization Planning Tool (AOPT).

Accounts for "risk to force" via penalties.

## **Our Contribution**

We model the problem as a two-player zero-sum game.

- The attacker selects a subset of the airfields to attack.
- The defender allocates the MOG (similarly to AOPT).

The literature on "attacker-defender" models (see e.g., [Alderson et al. 2011]) can be leveraged to compute prescribed MOG allocations.

# The Model: Sketch

$$\min_{Y} \max_{X,A,M} \left\{ X_{ts} - \sum_{i \in \mathcal{N}} q_i Y_i \left[ \sum_{j: (i,j) \in \mathcal{A}} X_{ij} + \sum_{j: (j,i) \in \mathcal{A}} X_{ji} \right] \right\}$$

#### **Attacker's Decision Variables:**

 $Y_i = \begin{cases} 1, & \text{if } i \in \mathcal{N} \text{ is attacked} \\ 0 & \text{otherwise} \end{cases}$ 

## Data:

- $\triangleright$   $\mathcal{N} =$ set of airfields (i.e., "nodes")
- A = set of feasible flight legs (i.e., "arcs")
- ▶  $q_i$  = penalty for flow through airfield  $i \in \mathbb{N}$ , if attacked
- Other data related to supply, cargo capacities of aircraft, . . .

## **Defender's Decision Variables**

- $X_{ij} = \text{amount of cargo sent on leg } (i, j)$
- $A_{ij}$  = number of aircraft sent on leg (i, j)
- $M_{ij}$  = amount of MOG moved from airfield *i* to *j*

**Constraints:** The  $M_{ij}$ 's constrain the  $A_{ij}$ 's, which constrain the  $X_{ij}$ 's.

# Full Formulation [Cooper 2023, page 9]

**ADAPT Formulation** 

$$\begin{split} \min_{I} \max_{X,A,M} & X_{ts} - \sum_{i \in \mathcal{N}} q_i Y_i \left[ \sum_{j: (i,j) \in \mathcal{R}} X_{ij} + \sum_{j: (j,i) \in \mathcal{R}} X_{ji} \right] \\ \text{subject to} & \sum_{j: (i,j) \in \mathcal{R}} X_{ij} - \sum_{j: (j,i) \in \mathcal{R}} X_{ji} = \begin{cases} X_{ts}, & \text{if } i = s \\ 0, & \text{if } i \notin \mathcal{N} \setminus \{s, t\} \\ -X_{ts} & \text{if } i = t \end{cases} \\ & X_{ij} \leq \begin{cases} cargocap \cdot A_{ij} & \text{if } (i, j) \in \mathcal{R} \\ supply & \text{if } (i, j) = (t, s) \end{cases} \\ & \sum_{j: (i,j) \in \mathcal{R}} M_{ij} \leq maxmogmoved \\ & - \sum_{j: (i,j) \in \mathcal{R}} M_{ij} + \sum_{j: (j,i) \in \mathcal{R}} M_{ji} \leq maxmog_i - startmog_i \quad \forall i \in \mathcal{N} \\ & \sum_{i \in \mathcal{N}} M_{ij} - \sum_{j: (i,j) \in \mathcal{R}} M_{ji} + \sum_{j: (j,i) \in \mathcal{R}} M_{ji} + \sum_{j: (j,i) \in \mathcal{R}} M_{ji} \leq startmog_i \quad \forall i \in \mathcal{N} \\ & \sum_{i \in \mathcal{N}} Y_i \leq maxattacks \\ & X_{ij}, A_{ij}, M_{ij} \geq 0 \quad \forall (i, j) \in \mathcal{R} \\ & Y_i \in \{0, 1\} \quad \forall i \in \mathcal{N} \end{split}$$



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Approved for public release. Distribution is unlimited.

## **Determining MOG Allocations**

#### **ADAPT Formulation**

$$\begin{split} \min_{Y \ X,A,M} & X_{is} - \sum_{i \in \mathcal{N}} q_i Y_i \left[ \sum_{j: (i,j) \in \mathcal{R}} X_{ij} + \sum_{j: (j,i) \in \mathcal{R}} X_{ji} \right] \\ \text{subject to} & \sum_{j: (i,j) \in \mathcal{R}} X_{ij} - \sum_{j: (j,i) \in \mathcal{R}} X_{ji} = \begin{cases} X_{is}, & \text{if } i = s \\ 0, & \text{if } i \notin \mathcal{N} \setminus \{s, t\} \\ -X_{is} & \text{if } i = t \end{cases} \\ & X_{ij} \leq \begin{cases} cargocap \cdot A_{ij} & \text{if } (i,j) \in \mathcal{R} \\ supply & \text{if } (i,j) = (t,s) \end{cases} \\ & \sum_{(i,j) \in \mathcal{R}} M_{ij} + \sum_{j: (j,i) \in \mathcal{R}} M_{ji} \leq \max mognoved \\ & - \sum_{j: (i,j) \in \mathcal{R}} M_{ij} - \sum_{j: (j,i) \in \mathcal{R}} M_{ji} + \sum_{j: (j,i) \in \mathcal{R}} M_{ji} \leq \max timog_i - startmog_i \quad \forall i \in \mathcal{N} \\ & \sum_{j: (i,j) \in \mathcal{R}} M_{ij} - \sum_{j: (j,i) \in \mathcal{R}} M_{ji} + \sum_{j: (j,i) \in \mathcal{R}} A_{ji} \\ & \sum_{i \in \mathcal{N}} Y_i \leq \max attacks \\ & X_{ij}, A_{ij}, M_{ij} \geq 0 \quad \forall (i,j) \in \mathcal{R} \\ & Y_i \in \{0, 1\} \quad \forall i \in \mathcal{N} \end{split}$$

## **Basic Idea**

- 1. Pick an initial attack Y.
- 2. Solve the defender's problem for *X*, *A*, *M*, given *Y*.
- 3. Try to find a better attack *Y*, given *X*, *A*, *M*.
- 4. Repeat . . .

More specifically (see [Cooper 2023, Section 3.2]),

- decompose the minimax problem into a "master" problem and a "defender subproblem", and
- iteratively add "Benders cuts" to the master problem.









# Take-Aways and Ongoing Work

## Take-Aways

- Formulated attacker-defender model for allocating airlift capacity in a contested environment.
- The model can be used to find "robust" MOG allocations.

#### **Ongoing Work**

- The model has been extended to supplying theater entry points with multiple cargo types, and penalties for unsatisfied demand.
- Part of an ongoing project sponsored by the Air Force Office of Scientific Research.